

Telnet Authentication: SRP

Status of this Memo

This document specifies an Internet standards track protocol for the Internet community, and requests discussion and suggestions for improvements. Please refer to the current edition of the "Internet Official Protocol Standards" (STD 1) for the standardization state and status of this protocol. Distribution of this memo is unlimited.

Copyright Notice

Copyright (C) The Internet Society (2000). All Rights Reserved.

Abstract

This document specifies an authentication scheme for the Telnet protocol under the framework described in [RFC2941], using the Secure Remote Password Protocol (SRP) authentication mechanism. The specific mechanism, SRP-SHA1, is described in [RFC2945].

1. Command Names and Codes

Authentication Types

SRP	5
-----	---

Suboption Commands

AUTH	0
REJECT	1
ACCEPT	2
CHALLENGE	3
RESPONSE	4
EXP	8
PARAMS	9

2. Command Meanings

IAC SB AUTHENTICATION IS <authentication-type-pair> AUTH IAC SE

This command indicates that the client has supplied the username and is ready to receive that user's field parameters. There is no authentication information to be sent to the remote side of the connection yet. This should only be sent after the IAC SB AUTHENTICATION NAME command has been issued. If the modifier byte (second byte of the authentication-type-pair) has any bits other than AUTH_WHO_MASK or AUTH_HOW_MASK set, both bytes are included in the session key hash described later. This ensures that the authentication type pair was correctly negotiated, while maintaining backward-compatibility with existing software.

IAC SB AUTHENTICATION REPLY <authentication-type-pair> PARAMS <values of modulus, generator, and salt> IAC SE

This command is used to pass the three parameter values used in the exponentiation to the client. These values are often called n , g , and s .

IAC SB AUTHENTICATION IS <authentication-type-pair> EXP <client's exponential residue> IAC SE

This command is used to pass the client's exponential residue, otherwise known as A , computed against the parameters exchanged earlier.

IAC SB AUTHENTICATION REPLY <authentication-type-pair> CHALLENGE <server's exponential residue> IAC SE

This command is used to pass the server's exponential residue, computed against the same parameters. This quantity is actually the sum of two residues, i.e. $g^x + g^b$. For details see [SRP] and [RFC2945].

IAC SB AUTHENTICATION IS <authentication-type-pair> RESPONSE <response from client> IAC SE

This command gives the server proof of the client's authenticity with a 160-bit (20 byte) response.

```
IAC SB AUTHENTICATION REPLY <authentication-type-pair> ACCEPT
<server's response> IAC SE
```

This command indicates that the authentication was successful. The server will construct its own proof of authenticity and include it as sub-option data.

```
IAC SB AUTHENTICATION REPLY <authentication-type-pair> REJECT
<optional reason for rejection> IAC SE
```

This command indicates that the authentication was not successful, and if there is any more data in the sub-option, it is an ASCII text message of the reason for the rejection.

For the PARAMS command, since three pieces of data are being transmitted, each parameter is preceded by a 16-bit (two byte) length specifier in network byte order. The EXP commands do not have a count in front of the data because there is only one piece of data in that suboption. The CHALLENGE, RESPONSE, and ACCEPT data also do not have a count because they are all fixed in size.

3. Implementation Rules

Currently, only AUTH_CLIENT_TO_SERVER mode is supported. Although the SRP protocol effectively performs implicit mutual authentication as a result of the two-way proofs, only the AUTH_HOW_ONE_WAY authentication mode is currently defined. The AUTH_HOW_MUTUAL setting is being reserved for an explicit mutual-authentication variant of the SRP protocol to be defined in future specifications.

All large number data sent in the arguments of the PARAMS and EXP commands must be in network byte order, i.e. most significant byte first. No padding is used.

The SRP-SHA1 mechanism, as described in [RFC2945] generates a 40-byte session key, which allows implementations to use different keys for incoming and outgoing traffic, increasing the security of the encrypted session. It is recommended that the Telnet ENCRYPT method, if it is used, be able to take advantage of the longer session keys.

4. Examples

User "tjw" may wish to log in on machine "foo". The client would send IAC SB AUTHENTICATION NAME "tjw" IAC SE IAC SB AUTHENTICATION IS SRP AUTH IAC SE. The server would look up the field and salt parameters for "tjw" from its password file and send them back to the client. Client and server would then exchange exponential residues and calculate their session keys (after the client prompted "tjw" for

his password). Then, the client would send the server its proof that it knows the session key. The server would either send back an ACCEPT or a REJECT. If the server accepts authentication, it also sends its own proof that it knows the session key to the client.

```
Client                                Server
IAC WILL AUTHENTICATION              IAC DO AUTHENTICATION

[ The server is now free to request authentication information.
  ]

                                      IAC SB AUTHENTICATION SEND
                                      SRP CLIENT|ONE_WAY|
                                      ENCRYPT_USING_TELOPT
                                      SRP CLIENT|ONE_WAY
                                      IAC SE
```

[The server has requested SRP authentication. It has indicated a preference for ENCRYPT_USING_TELOPT, which requires the Telnet ENCRYPT option to be negotiated once authentication succeeds. If the client does not support this, the server is willing to fall back to an encryption-optional mode.

The client will now respond with the name of the user that it wants to log in as.]

```
IAC SB AUTHENTICATION NAME
"tjw" IAC SE
IAC SB AUTHENTICATION IS
SRP CLIENT|ONE_WAY|ENCRYPT_USING_TELOPT AUTH
IAC SE
```

[The server looks up the appropriate information for "tjw" and sends back the parameters in a PARAMS command. The parameters consist of the values N, g, and s, each preceded with a two-byte size parameter.]

```
IAC SB AUTHENTICATION REPLY
SRP CLIENT|ONE_WAY|
ENCRYPT_USING_TELOPT PARAMS
ss ss nn nn nn nn ...
ss ss gg gg gg gg ...
ss ss tt tt tt tt ...
IAC SE
```

[Both sides send their exponential residues. The client sends its value A and the server sends its value B. In SRP, the CHALLENGE message may be computed but not sent before the EXP command.]

```
IAC SB AUTHENTICATION IS
SRP CLIENT|ONE_WAY|ENCRYPT_USING_TELOPT EXP
aa aa aa aa aa aa aa ...
IAC SE
```

```
IAC SB AUTHENTICATION REPLY
SRP CLIENT|ONE_WAY|
ENCRYPT_USING_TELOPT CHALLENGE
bb bb bb bb bb bb bb ...
IAC SE
```

[The client sends its response to the server. This is the message M in the SRP protocol, which proves possession of the session key by the client.

Since ENCRYPT_USING_TELOPT is specified, the two octets of the authentication-type-pair are appended to the session key K before the hash for M is computed. If the client and server had agreed upon a mode without the encryption flag set, nothing would be appended to K.

Both this message and the server's response are as long as the output of the hash; the length is 20 bytes for SHA-1.]

```
IAC SB AUTHENTICATION IS
SRP CLIENT|ONE_WAY|ENCRYPT_USING_TELOPT RESPONSE
xx xx xx xx xx xx xx ...
IAC SE
```

[The server accepts the response and sends its own proof.]

```
IAC SB AUTHENTICATION REPLY
SRP CLIENT|ONE_WAY|
ENCRYPT_USING_TELOPT ACCEPT
yy yy yy yy yy yy yy ...
IAC SE
```

5. Security Considerations

The ability to negotiate a common authentication mechanism between client and server is a feature of the authentication option that should be used with caution. When the negotiation is performed, no authentication has yet occurred. Therefore, each system has no way of knowing whether or not it is talking to the system it intends. An intruder could attempt to negotiate the use of an authentication system which is either weak, or already compromised by the intruder.

Since SRP relies on the security of the underlying public-key cryptosystem, the modulus "n" should be large enough to resist brute-force attack. A length of at least 1024 bits is recommended, and implementations should reject attempts to use moduli that are shorter than 512 bits, or attempts to use invalid moduli and generator parameters (non-safe-prime "n" or non-primitive "g").

6. IANA Considerations

The authentication type SRP and its associated suboption values are registered with IANA. Any suboption values used to extend the protocol as described in this document must be registered with IANA before use. IANA is instructed not to issue new suboption values without submission of documentation of their use.

7. References

- [RFC2941] Ts'o, T. and J. Altman, "Telnet Authentication Option", RFC 2941, September 2000.
- [SRP] T. Wu, "The Secure Remote Password Protocol", In Proceedings of the 1998 ISOC Network and Distributed System Security Symposium, San Diego, CA, pp. 97-111.
- [RFC2945] Wu, T., "The SRP Authentication and Key Exchange System", RFC 2945, September 2000.

8. Author's Address

Thomas Wu
Stanford University
Stanford, CA 94305

EMail: tjw@cs.Stanford.EDU

9. Full Copyright Statement

Copyright (C) The Internet Society (2000). All Rights Reserved.

This document and translations of it may be copied and furnished to others, and derivative works that comment on or otherwise explain it or assist in its implementation may be prepared, copied, published and distributed, in whole or in part, without restriction of any kind, provided that the above copyright notice and this paragraph are included on all such copies and derivative works. However, this document itself may not be modified in any way, such as by removing the copyright notice or references to the Internet Society or other Internet organizations, except as needed for the purpose of developing Internet standards in which case the procedures for copyrights defined in the Internet Standards process must be followed, or as required to translate it into languages other than English.

The limited permissions granted above are perpetual and will not be revoked by the Internet Society or its successors or assigns.

This document and the information contained herein is provided on an "AS IS" basis and THE INTERNET SOCIETY AND THE INTERNET ENGINEERING TASK FORCE DISCLAIMS ALL WARRANTIES, EXPRESS OR IMPLIED, INCLUDING BUT NOT LIMITED TO ANY WARRANTY THAT THE USE OF THE INFORMATION HEREIN WILL NOT INFRINGE ANY RIGHTS OR ANY IMPLIED WARRANTIES OF MERCHANTABILITY OR FITNESS FOR A PARTICULAR PURPOSE.

Acknowledgement

Funding for the RFC Editor function is currently provided by the Internet Society.

