                Stateless IP/ICMP Translation Algorithm (SIIT)

Status of this Memo

   This document specifies an Internet standards track protocol for the
   Internet community, and requests discussion and suggestions for
   improvements. Please refer to the current edition of the "Internet
   Official Protocol Standards" (STD 1) for the standardization state
   and status of this protocol. Distribution of this memo is unlimited.

Abstract

   This document specifies a transition mechanism algorithm in addition
   to the mechanisms already specified in [TRANS-MECH]. The algorithm
   translates between IPv4 and IPv6 packet headers (including ICMP
   headers) in separate translator "boxes" in the network without
   requiring any per-connection state in those "boxes". This new
   algorithm can be used as part of a solution that allows IPv6 hosts,
   which do not have a permanently assigned IPv4 addresses, to
   communicate with IPv4-only hosts. The document neither specifies
   address assignment nor routing to and from the IPv6 hosts when they
   communicate with the IPv4-only hosts.

Table of Contents

1. Introduction and Motivation

   The transition mechanisms specified in [TRANS-MECH] handle the case
   of dual IPv4/IPv6 hosts interoperating with both dual hosts and
   IPv4-only hosts, which is needed early in the transition to IPv6.
   The dual hosts are assigned both an IPv4 and one or more IPv6
   addresses. As the number of available globally unique IPv4 addresses
   becomes smaller and smaller as the Internet grows there will be a
   desire to take advantage of the large IPv6 address and not require
   that every new Internet node have a permanently assigned IPv4
   address.

   There are several different scenarios where there might be IPv6-only
   hosts that need to communicate with IPv4-only hosts. These IPv6
   hosts might be IPv4-capable, i.e. include an IPv4 implementation but
   not be assigned an IPv4 address, or they might not even include an
   IPv4 implementation.

   - A completely new network with new devices that all support IPv6.
      In this case it might be beneficial to not have to configure the
      routers within the new network to route IPv4 since none of the

hosts in the new network are configured with IPv4 addresses. But
these new IPv6 devices might occasionally need to communicate with
some IPv4 nodes out on the Internet.

- An existing network where a large number of IPv6 devices are
   added. The IPv6 devices might have both an IPv4 and an IPv6
   protocol stack but there is not enough global IPv4 address space
   to give each one of them a permanent IPv4 address. In this case
   it is more likely that the routers in the network already route
   IPv4 and are upgraded to dual routers.

However, there are other potential solutions in this area:

- If there is no IPv4 routing inside the network i.e., the cloud
   that contains the new devices, some possible solutions are to
   either use the translators specified in this document at the
   boundary of the cloud, or to use Application Layer Gateways (ALG)
   on dual nodes at the cloud's boundary. The ALG solution is less
   flexible in that it is application protocol specific and it is
   also less robust since an ALG box is likely to be a single point
   of failure for a connection using that box.

- Otherwise, if IPv4 routing is supported inside the cloud and the
   implementations support both IPv6 and IPv4 it might suffice to
   have a mechanism for allocating a temporary address IPv4 and use
   IPv4 end to end when communicating with IPv4-only nodes. However,
   it would seem that such a solution would require the pool of
   temporary IPv4 addresses to be partitioned across all the subnets
   in the cloud which would either require a larger pool of IPv4
   addresses or result in cases where communication would fail due to
   no available IPv4 address for the node's subnet.

This document specifies an algorithm that is one of the components
needed to make IPv6-only nodes interoperate with IPv4-only nodes.
Other components, not specified in this document, are a mechanism for
the IPv6-only node to somehow acquire a temporary IPv4 address, and a
mechanism for providing routing (perhaps using tunneling) to and from
the temporary IPv4 address assigned to the node.

The temporary IPv4 address will be used as an IPv4-translated IPv6
address and the packets will travel through a stateless IP/ICMP
translator that will translate the packet headers between IPv4 and
IPv6 and translate the addresses in those headers between IPv4
addresses on one side and IPv4-translated or IPv4-mapped IPv6
addresses on the other side.

This specification does not cover how an IPv6 node can acquire a
temporary IPv4 address and how such a temporary address be registered
in the DNS. The DHCP protocol, perhaps with some extensions, could
probably be used to acquire temporary addresses with short leases but
that is outside the scope of this document. Also, the mechanism for
routing this IPv4-translated IPv6 address in the site is not
specified in this document.

The figures below show how the Stateless IP/ICMP Translation
algorithm (SIIT) can be used initially for small networks (e.g., a
single subnet) and later for a site which has IPv6-only hosts in a
dual IPv4/IPv6 network. This use assumes a mechanism for the IPv6
nodes to acquire a temporary address from the pool of IPv4 addresses.
Note that SIIT is not likely to be useful later during transition
when most of the Internet is IPv6 and there are only small islands of
IPv4 nodes, since such use would either require the IPv6 nodes to
acquire temporary IPv4 addresses from a "distant" SIIT box operated
by a different administration, or require that the IPv6 routing
contain routes for IPv6-mapped addresses. (The latter is known to be
a very bad idea due to the size of the IPv4 routing table that would
potentially be injected into IPv6 routing in the form of IPv4-mapped
addresses.)

```
                                     _____
                                    /             \
      [IPv6 Host]---[SIIT]---------<  IPv4 network>--[IPv4 Host]
                      |             _____/
              (pool of IPv4 addresses)

      IPv4-translatable ->            IPv4->IPv4 addresser
      IPv4-mapped
```
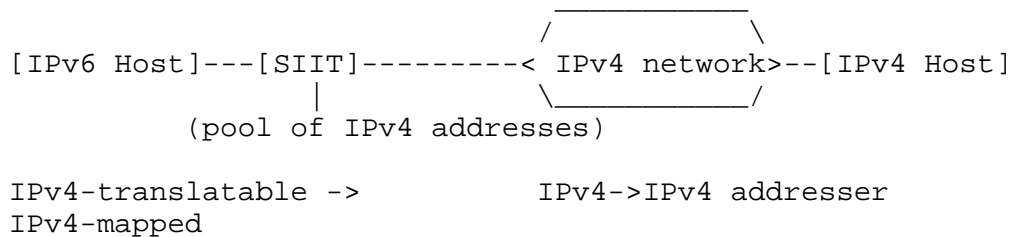
         Figure 1. Using SIIT for a single IPv6-only subnet.

```
              _____                 _____
             /             \               /             \
      [IPv6 Host]--< Dual network>--[SIIT]--< IPv4 network>--[IPv4 Host]
             _____/      |       _____/
                       (pool of IPv4 addresses)

      IPv4-translatable ->                   IPv4->IPv4 addresser
      IPv4-mapped
```
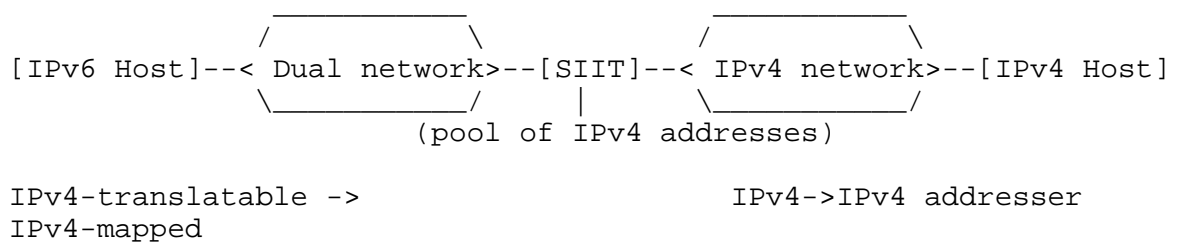
   Figure 2. Using SIIT for an IPv6-only or dual cloud (e.g. a site)
        which contains some IPv6-only hosts as well as IPv4 hosts.

The protocol translators are assumed to fit around some piece of
topology that includes some IPv6-only nodes and that may also include
IPv4 nodes as well as dual nodes. There has to be a translator on
each path used by routing the "translatable" packets in and out of
this cloud to ensure that such packets always get translated. This
does not require a translator at every physical connection between
the cloud and the rest of the Internet since the routing can be used
to deliver the packets to the translator.

The IPv6-only node communicating with an IPv4 node through a
translator will see an IPv4-mapped address for the peer and use an
IPv4-translatable address for its local address for that
communication. When the IPv6-only node sends packets the IPv4-mapped
address indicates that the translator needs to translate the packets.
When the IPv4 node sends packets those will translated to have the
IPv4-translatable address as a destination; it is not possible to use
an IPv4-mapped or an IPv4-compatible address as a destination since
that would either route the packet back to the translator (for the
IPv4-mapped address) or make the packet be encapsulated in IPv4 (for
the IPv4-compatible address). Thus this specification introduces the
new notion of an IPv4-translatable address.

1.1. Applicability and Limitations

The use of this translation algorithm assumes that the IPv6 network
is somehow well connected i.e. when an IPv6 node wants to communicate
with another IPv6 node there is an IPv6 path between them. Various
tunneling schemes exist that can provide such a path, but those
mechanisms and their use is outside the scope of this document.

The IPv6 protocol [IPv6] has been designed so that the TCP and UDP
pseudo-header checksums are not affected by the translations
specified in this document, thus the translator does not need to
modify normal TCP and UDP headers. The only exceptions are
unfragmented IPv4 UDP packets which need to have a UDP checksum
computed since a pseudo-header checksum is required for UDP in IPv6.
Also, ICMPv6 include a pseudo-header checksum but it is not present
in ICMPv4 thus the checksum in ICMP messages need to be modified by
the translator. In addition, ICMP error messages contain an IP
header as part of the payload thus the translator need to rewrite
those parts of the packets to make the receiver be able to understand
the included IP header. However, all of the translator's operations,
including path MTU discovery, are stateless in the sense that the
translator operates independently on each packet and does not retain
any state from one packet to another. This allows redundant
translator boxes without any coordination and a given TCP connection
can have the two directions of packets go through different
translator boxes.

The translating function as specified in this document does not
translate any IPv4 options and it does not translate IPv6 routing
headers, hop-by-hop extension headers, or destination options
headers. It could be possible to define a translation between source
routing in IPv4 and IPv6. However such a translation would not be
semantically correct due to the slight differences between the IPv4
and IPv6 source routing. Also, the usefulness of source routing when
going through a header translator might be limited since all the
IPv6-only routers would need to have an IPv4-translated IPv6 address
since the IPv4-only node will send a source route option containing
only IPv4 addresses.

At first sight it might appear that the IPsec functionality [IPv6-SA,
IPv6-ESP, IPv6-AH] can not be carried across the translator.
However, since the translator does not modify any headers above the
logical IP layer (IP headers, IPv6 fragment headers, and ICMP
messages) packets encrypted using ESP in Transport-mode can be
carried through the translator. [Note that this assumes that the key
management can operate between the IPv6-only node and the IPv4-only
node.] The AH computation covers parts of the IPv4 header fields
such as IP addresses, and the identification field (fields that are
either immutable or predictable by the sender) [IPv6-AUTH]. While
the SIIT algorithm is specified so that those IPv4 fields can be
predicted by the IPv6 sender it is not possible for the IPv6 receiver
to determine the value of the IPv4 Identification field in packets
sent by the IPv4 node. Thus as the translation algorithm is
specified in this document it is not possible to use end-to-end AH
through the translator.

For ESP Tunnel-mode to work through the translator the IPv6 node
would have to be able to both parse and generate "inner" IPv4 headers
since the inner IP will be encrypted together with the transport
protocol.

Thus in practise, only ESP transport mode is relatively easy to make
work through a translator.

IPv4 multicast addresses can not be mapped to IPv6 multicast
addresses. For instance, ::ffff:224.1.2.3 is an IPv4 mapped IPv6
address with a class D address, however it is not an IPv6 multicast
address. While the IP/ICMP header translation aspect of this memo in
theory works for multicast packets this address mapping limitation
makes it impossible to apply the techniques in this memo for
multicast traffic.

1.2. Assumptions

   The IPv6 nodes using the translator must have an IPv4-translated IPv6
   address while it is communicating with IPv4-only nodes.

   The use of the algorithm assumes that there is an IPv4 address pool
   used to generate IPv4-translated addresses. Routing needs to be able
   to route any IPv4 packets, whether generated "outside" or "inside"
   the translator, destined to addresses in this pool towards the
   translator. This implies that the address pool can not be assigned
   to subnets but must be separated from the IPv4 subnets used on the
   "inside" of the translator.

   Fragmented IPv4 UDP packets that do not contain a UDP checksum (i.e.
   the UDP checksum field is zero) are not of significant use over
   wide-areas in the Internet and will not be translated by the
   translator. An informal trace [MILLER] in the backbone showed that
   out of 34,984,468 IP packets there were 769 fragmented UDP packets
   with a zero checksum. However, all of them were due to malicious or
   broken behavior; a port scan and first fragments of IP packets that
   are not a multiple of 8 bytes.

1.3. Impact Outside the Network Layer

   The potential existence of stateless IP/ICMP translators is already
   taken care of from a protocol perspective in [IPv6]. However, an
   IPv6 node that wants to be able to use translators needs some
   additional logic in the network layer.

   The network layer in an IPv6-only node, when presented by the
   application with either an IPv4 destination address or an IPv4-mapped
   IPv6 destination address, is likely to drop the packet and return
   some error message to the application. In order to take advantage of
   translators such a node should instead send an IPv6 packet where the
   destination address is the IPv4-mapped address and the source address
   is the node's temporarily assigned IPv4-translated address. If the
   node does not have a temporarily assigned IPv4-translated address it
   should acquire one using mechanisms that are not discussed in this
   document.

   Note that the above also applies to a dual IPv4/IPv6 implementation
   node which is not configured with any IPv4 address.

   There are no extra changes needed to applications to operate through
   a translator beyond what applications already need to do to operate
   on a dual node. The applications that have been modified to work on
   a dual node already have the mechanisms to determine whether they are
   communicating with an IPv4 or an IPv6 peer. Thus if the applications

need to modify their behavior depending on the type of the peer, such
as ftp determining whether to fallback to using the PORT/PASV command
when EPRT/EPSV fails (as specified in [FTPEXT]), they already need to
do that when running on dual nodes and the presense of translators
does not add anything. For example, when using the socket API
[BSDAPI] the applications know that the peer is IPv6 if they get an
AF_INET6 address from the name service and the address is not an
IPv4-mapped address (i.e., IN6_IS_ADDR_V4MAPPED returns false). If
this is not the case, i.e., the address is AF_INET or an IPv4-mapped
IPv6 address, the peer is IPv4.

One way of viewing the translator, which might help clarify why
applications do not need to know that a translator is used, is to
look at the information that is passed from the transport layer to
the network layer. If the transport passes down an IPv4 address
(whether or not is in the IPv4-mapped encoding) this means that at
some point there will be IPv4 packets generated. In a dual node the
generation of the IPv4 packets takes place in the sending node. In
an IPv6-only node conceptually the only difference is that the IPv4
packet is generated by the translator - all the information that the
transport layer passed to the network layer will be conveyed to the
translator in some form. That form just "happens" to be in the form
of an IPv6 header.

2. Terminology

   This documents uses the terminology defined in [IPv6] and
   [TRANS-MECH] with these clarifications:

        IPv4 capable node:
                A node which has an IPv4 protocol stack.
                In order for the stack to be usable the node must be
                assigned one or more IPv4 addresses.

        IPv4 enabled node:
                A node which has an IPv4 protocol stack
                and is assigned one or more IPv4 addresses. Both
                IPv4-only and IPv6/IPv4 nodes are IPv4 enabled.

        IPv6 capable node:
                A node which has an IPv6 protocol stack.
                In order for the stack to be usable the node must be
                assigned one or more IPv6 addresses.

        IPv6 enabled node:
                A node which has an IPv6 protocol stack
                and is assigned one or more IPv6 addresses. Both
                IPv6-only and IPv6/IPv4 nodes are IPv6 enabled.

2.1. Addresses

   In addition to the forms of addresses defined in [ADDR-ARCH] this
   document also introduces the new form of IPv4-translated address.
   This is needed to avoid using IPv4-compatible addresses outside the
   intended use of automatic tunneling. Thus the address forms are:

        IPv4-mapped:
                An address of the form 0::ffff:a.b.c.d which refers
                to a node that is not IPv6-capable. In addition to
                its use in the API this protocol uses IPv4-mapped
                addresses in IPv6 packets to refer to an IPv4 node.

        IPv4-compatible:
                An address of the form 0::0:a.b.c.d which refers to
                an IPv6/IPv4 node that supports automatic tunneling.
                Such addresses are not used in this protocol.

        IPv4-translated:
                An address of the form 0::ffff:0:a.b.c.d which refers
                to an IPv6-enabled node. Note that the prefix
                0::ffff:0:0:0/96 is chosen to checksum to zero to
                avoid any changes to the transport protocol's pseudo
                header checksum.

2.2. Requirements

   The keywords MUST, MUST NOT, REQUIRED, SHALL, SHALL NOT, SHOULD,
   SHOULD NOT, RECOMMENDED, MAY, and OPTIONAL, when they appear in this
   document, are to be interpreted as described in [KEYWORDS].

3. Translating from IPv4 to IPv6

   When an IPv4-to-IPv6 translator receives an IPv4 datagram addressed
   to a destination that lies outside of the attached IPv4 island, it
   translates the IPv4 header of that packet into an IPv6 header. It
   then forwards the packet based on the IPv6 destination address. The
   original IPv4 header on the packet is removed and replaced by an IPv6
   header. Except for ICMP packets the transport layer header and data
   portion of the packet are left unchanged.

```
       +-------------+                      +-------------+
       |    IPv4     |                      |    IPv6     |
       |   Header    |                      |   Header    |
       +-------------+                      +-------------+
       |  Transport  |                      |  Fragment   |
       |    Layer    |       ===>           |   Header    |
       |   Header    |                      | (not always)|
       +-------------+                      +-------------+
       |             |                      |  Transport  |
       ~    Data     ~                      |    Layer    |
       |             |                      |   Header    |
       +-------------+                      +-------------+
                                            |             |
                                            ~    Data     ~
                                            |             |
                                            +-------------+
```
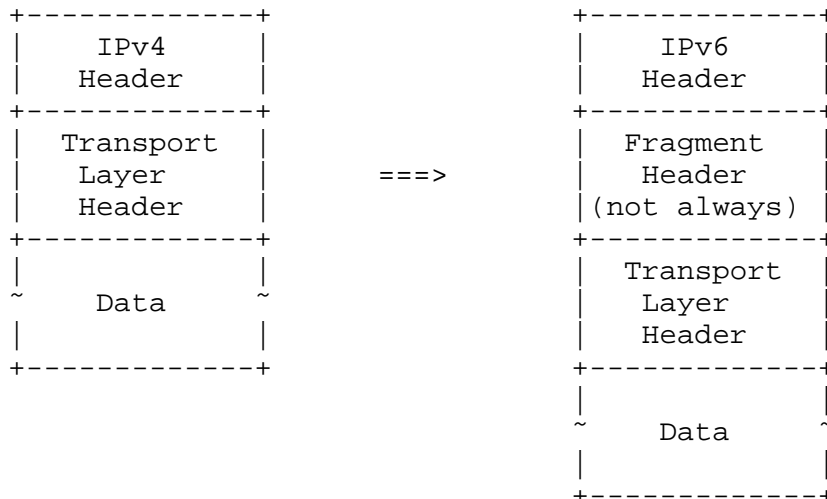
                      IPv4-to-IPv6 Translation

   One of the differences between IPv4 and IPv6 is that in IPv6 path MTU
   discovery is mandatory but it is optional in IPv4. This implies that
   IPv6 routers will never fragment a packet - only the sender can do
   fragmentation.

   When the IPv4 node performs path MTU discovery (by setting the DF bit
   in the header) the path MTU discovery can operate end-to-end i.e.
   across the translator. In this case either IPv4 or IPv6 routers
   might send back ICMP "packet too big" messages to the sender. When
   these ICMP errors are sent by the IPv6 routers they will pass through
   a translator which will translate the ICMP error to a form that the
   IPv4 sender can understand. In this case an IPv6 fragment header is
   only included if the IPv4 packet is already fragmented.

   However, when the IPv4 sender does not perform path MTU discovery the
   translator has to ensure that the packet does not exceed the path MTU
   on the IPv6 side. This is done by fragmenting the IPv4 packet so
   that it fits in 1280 byte IPv6 packet since IPv6 guarantees that 1280
   byte packets never need to be fragmented. Also, when the IPv4 sender
   does not perform path MTU discovery the translator MUST always
   include an IPv6 fragment header to indicate that the sender allows
   fragmentation. That is needed should the packet pass through an
   IPv6-to-IPv4 translator.

   The above rules ensure that when packets are fragmented either by the
   sender or by IPv4 routers that the low-order 16 bits of the fragment
   identification is carried end-end to ensure that packets are
   correctly reassembled. In addition, the rules use the presence of an

IPv6 fragment header to indicate that the sender might not be using
path MTU discovery i.e. the packet should not have the DF flag set
should it later be translated back to IPv4.

Other than the special rules for handling fragments and path MTU
discovery the actual translation of the packet header consists of a
simple mapping as defined below. Note that ICMP packets require
special handling in order to translate the content of ICMP error
message and also to add the ICMP pseudo-header checksum.

3.1. Translating IPv4 Headers into IPv6 Headers

If the DF flag is not set and the IPv4 packet will result in an IPv6
packet larger than 1280 bytes the IPv4 packet MUST be fragmented
prior to translating it. Since IPv4 packets with DF not set will
always result in a fragment header being added to the packet the IPv4
packets must be fragmented so that their length, excluding the IPv4
header, is at most 1232 bytes (1280 minus 40 for the IPv6 header and
8 for the Fragment header). The resulting fragments are then
translated independently using the logic described below.

If the DF bit is set and the packet is not a fragment (i.e., the MF
flag is not set and the Fragment Offset is zero) then there is no
need to add a fragment header to the packet. The IPv6 header fields
are set as follows:

        Version:
                6

        Traffic Class:
                By default, copied from IP Type Of Service and
                Precedence field (all 8 bits are copied). According
                to [DIFFSERV] the semantics of the bits are identical
                in IPv4 and IPv6. However, in some IPv4 environments
                these fields might be used with the old semantics of
                "Type Of Service and Precedence". An implementation
                of a translator SHOULD provide the ability to ignore
                the IPv4 "TOS" and always set the IPv6 traffic class
                to zero.

        Flow Label:
                0 (all zero bits)

        Payload Length:
                Total length value from IPv4 header, minus the size
                of the IPv4 header and IPv4 options, if present.

Next Header:
        Protocol field copied from IPv4 header

Hop Limit:
        TTL value copied from IPv4 header. Since the
        translator is a router, as part of forwarding the
        packet it needs to decrement either the IPv4 TTL
        (before the translation) or the IPv6 Hop Limit (after
        the translation). As part of decrementing the TTL or
        Hop Limit the translator (as any router) needs to
        check for zero and send the ICMPv4 or ICMPv6 "ttl
        exceeded" error.

Source Address:
        The low-order 32 bits is the IPv4 source address.
        The high-order 96 bits is the IPv4-mapped prefix
        (::ffff:0:0/96)

Destination Address:
        The low-order 32 bits is the IPv4 destination
        address. The high-order 96 bits is the IPv4-
        translated prefix (0::ffff:0:0:0/96)

If IPv4 options are present in the IPv4 packet, they are ignored
i.e., there is no attempt to translate them. However, if an
unexpired source route option is present then the packet MUST instead
be discarded, and an ICMPv4 "destination unreachable/source route
failed" (Type 3/Code 5) error message SHOULD be returned to the
sender.

If there is need to add a fragment header (the DF bit is not set or
the packet is a fragment) the header fields are set as above with the
following exceptions:

   IPv6 fields:

      Payload Length:
              Total length value from IPv4 header, plus 8 for the
              fragment header, minus the size of the IPv4 header
              and IPv4 options, if present.

      Next Header:
              Fragment Header (44).

   Fragment header fields:

      Next Header:
              Protocol field copied from IPv4 header.

Fragment Offset:
        Fragment Offset copied from the IPv4 header.

M flag:
        More Fragments bit copied from the IPv4 header.

Identification:
        The low-order 16 bits copied from the Identification
        field in the IPv4 header. The high-order 16 bits set
        to zero.

3.2. Translating UDP over IPv4

If a UDP packet has a zero UDP checksum then a valid checksum must be
calculated in order to translate the packet. A stateless translator
can not do this for fragmented packets but [MILLER] indicates that
fragmented UDP packets with a zero checksum appear to only be used
for malicious purposes. Thus this is not believed to be a noticeable
limitation.

When a translator receives the first fragment of a fragmented UDP
IPv4 packet and the checksum field is zero the translator SHOULD drop
the packet and generate a system management event specifying at least
the IP addresses and port numbers in the packet. When it receives
fragments other than the first it SHOULD silently drop the packet,
since there is no port information to log.

When a translator receives an unfragmented UDP IPv4 packet and the
checksum field is zero the translator MUST compute the missing UDP
checksum as part of translating the packet. Also, the translator
SHOULD maintain a counter of how many UDP checksums are generated in
this manner.

3.3. Translating ICMPv4 Headers into ICMPv6 Headers

All ICMP messages that are to be translated require that the ICMP
checksum field be updated as part of the translation since ICMPv6,
unlike ICMPv4, has a pseudo-header checksum just like UDP and TCP.

In addition all ICMP packets need to have the Type value translated
and for ICMP error messages the included IP header also needs
translation.

The actions needed to translate various ICMPv4 messages are:

   ICMPv4 query messages:

      Echo and Echo Reply (Type 8 and Type 0)
         Adjust the type to 128 and 129, respectively, and adjust the
         ICMP checksum both to take the type change into account and
         to include the ICMPv6 pseudo-header.

      Information Request/Reply (Type 15 and Type 16)
         Obsoleted in ICMPv4. Silently drop.

      Timestamp and Timestamp Reply (Type 13 and Type 14)
         Obsoleted in ICMPv6. Silently drop.

      Address Mask Request/Reply (Type 17 and Type 18)
         Obsoleted in ICMPv6. Silently drop.

      ICMP Router Advertisement (Type 9)
         Single hop message. Silently drop.

      ICMP Router Solicitation (Type 10)
         Single hop message. Silently drop.

      Unknown ICMPv4 types
         Silently drop.

   IGMP messages:

         While the MLD messages [MLD] are the logical IPv6
         counterparts for the IPv4 IGMP messages all the "normal" IGMP
         messages are single-hop messages and should be silently
         dropped by the translator. Other IGMP messages might be used
         by multicast routing protocols and, since it would be a
         configuration error to try to have router adjacencies across
         IPv4/IPv6 translators those packets should also be silently
         dropped.

   ICMPv4 error messages:

     Destination Unreachable (Type 3)
        For all that are not explicitly listed below set the Type to
        1.

        Translate the code field as follows:
           Code 0, 1 (net, host unreachable):
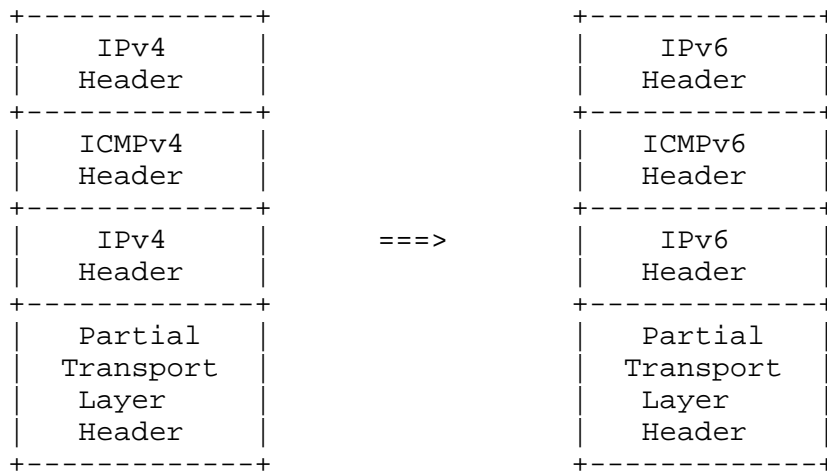   Set Code to 0 (no route to destination).

        Code 2 (protocol unreachable):
             Translate to an ICMPv6 Parameter Problem (Type 4,
             Code 1) and make the Pointer point to the IPv6 Next
             Header field.

        Code 3 (port unreachable):
             Set Code to 4 (port unreachable).

        Code 4 (fragmentation needed and DF set):
             Translate to an ICMPv6 Packet Too Big message (Type
             2) with code 0. The MTU field needs to be adjusted
             for the difference between the IPv4 and IPv6 header
             sizes. Note that if the IPv4 router did not set
             the MTU field i.e. the router does not implement
             [PMTUv4], then the translator must use the plateau
             values specified in [PMTUv4] to determine a likely
             path MTU and include that path MTU in the ICMPv6
             packet. (Use the greatest plateau value that is
             less than the returned Total Length field.)

        Code 5 (source route failed):
             Set Code to 0 (no route to destination). Note that
             this error is unlikely since source routes are not
             translated.

        Code 6,7:
             Set Code to 0 (no route to destination).

        Code 8:
             Set Code to 0 (no route to destination).

        Code 9, 10 (communication with destination host
        administratively prohibited):
             Set Code to 1 (communication with destination
             administratively prohibited)

        Code 11, 12:
             Set Code to 0 (no route to destination).

   Redirect (Type 5)
      Single hop message. Silently drop.

   Source Quench (Type 4)
      Obsoleted in ICMPv6. Silently drop.

   Time Exceeded (Type 11)
      Set the Type field to 3. The Code field is unchanged.

            Parameter Problem (Type 12)
                Set the Type field to 4. The Pointer needs to be updated to
                point to the corresponding field in the translated include
                IP header.

3.4. Translating ICMPv4 Error Messages into ICMPv6

    There are some differences between the IPv4 and the IPv6 ICMP error
    message formats as detailed above. In addition, the ICMP error
    messages contain the IP header for the packet in error which needs to
    be translated just like a normal IP header. The translation of this
    "packet in error" is likely to change the length of the datagram thus
    the Payload Length field in the outer IPv6 header might need to be
    updated.

```
            +-------------+                        +-------------+
            |    IPv4     |                        |    IPv6     |
            |   Header    |                        |   Header    |
            +-------------+                        +-------------+
            |   ICMPv4    |                        |   ICMPv6    |
            |   Header    |                        |   Header    |
            +-------------+                        +-------------+
            |    IPv4     |        ===>            |    IPv6     |
            |   Header    |                        |   Header    |
            +-------------+                        +-------------+
            |   Partial   |                        |   Partial   |
            | Transport   |                        | Transport   |
            |   Layer     |                        |   Layer     |
            |   Header    |                        |   Header    |
            +-------------+                        +-------------+
```

                IPv4-to-IPv6 ICMP Error Translation

    The translation of the inner IP header can be done by recursively
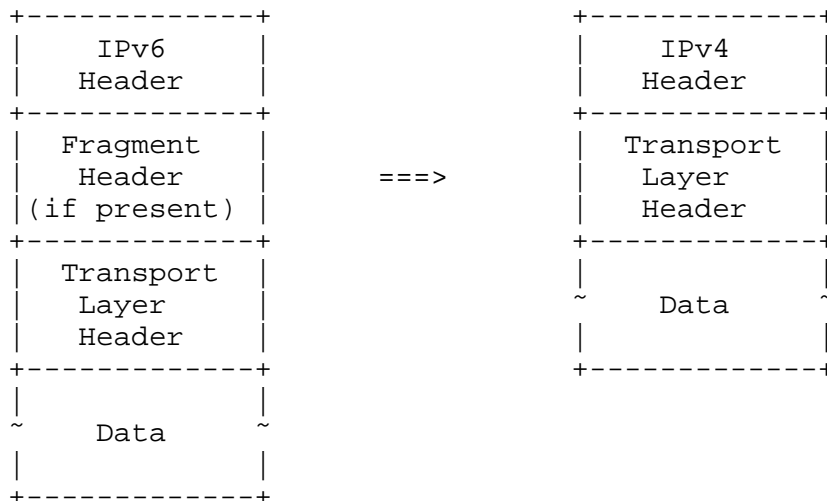   invoking the function that translated the outer IP headers.

3.5. Knowing when to Translate

    The translator is assumed to know the pool(s) of IPv4 address that
    are used to represent the internal IPv6-only nodes. Thus if the IPv4
    destination field contains an address that falls in these configured
    sets of prefixes the packet needs to be translated to IPv6.

4. Translating from IPv6 to IPv4

   When an IPv6-to-IPv4 translator receives an IPv6 datagram addressed
   to an IPv4-mapped IPv6 address, it translates the IPv6 header of that
   packet into an IPv4 header. It then forwards the packet based on the
   IPv4 destination address. The original IPv6 header on the packet is
   removed and replaced by an IPv4 header. Except for ICMP packets the
   transport layer header and data portion of the packet are left
   unchanged.

```
           +------------+                 +------------+
           |    IPv6    |                 |    IPv4    |
           |   Header   |                 |   Header   |
           +------------+                 +------------+
           |  Fragment  |                 | Transport  |
           |   Header   |      ===>       |   Layer    |
           |(if present)|                 |   Header   |
           +------------+                 +------------+
           | Transport  |                 |            |
           |   Layer    |                 ~    Data    ~
           |   Header   |                 |            |
           +------------+                 +------------+
           |            |
           ~    Data    ~
           |            |
           +------------+
```

                    IPv6-to-IPv4 Translation

   There are some differences between IPv6 and IPv4 in the area of
   fragmentation and the minimum link MTU that effect the translation.
   An IPv6 link has to have an MTU of 1280 bytes or greater. The
   corresponding limit for IPv4 is 68 bytes. Thus, unless there were
   special measures, it would not be possible to do end-to-end path MTU
   discovery when the path includes an IPv6-to-IPv4 translator since the
   IPv6 node might receive ICMP "packet too big" messages originated by
   an IPv4 router that report an MTU less than 1280. However, [IPv6]
   requires that IPv6 nodes handle such an ICMP "packet too big" message
   by reducing the path MTU to 1280 and including an IPv6 fragment
   header with each packet. This allows end-to-end path MTU discovery
   across the translator as long as the path MTU is 1280 bytes or
   greater. When the path MTU drops below the 1280 limit the IPv6
   sender will originate 1280 byte packets that will be fragmented by
   IPv4 routers along the path after being translated to IPv4.

   The only drawback with this scheme is that it is not possible to use
   PMTU to do optimal UDP fragmentation (as opposed to completely
   avoiding fragmentation) at sender since the presence of an IPv6

Fragment header is interpreted that is it OK to fragment the packet
on the IPv4 side. Thus if a UDP application wants to send large
packets independent of the PMTU, the sender will only be able to
determine the path MTU on the IPv6 side of the translator. If the
path MTU on the IPv4 side of the translator is smaller then the IPv6
sender will not receive any ICMP "too big" errors and can not adjust
the size fragments it is sending.

Other than the special rules for handling fragments and path MTU
discovery the actual translation of the packet header consists of a
simple mapping as defined below. Note that ICMP packets require
special handling in order to translate the content of ICMP error
message and also to add the ICMP pseudo-header checksum.

4.1. Translating IPv6 Headers into IPv4 Headers

If there is no IPv6 Fragment header the IPv4 header fields are set as
follows:

> Version:
>> 4
>
> Internet Header Length:
>> 5 (no IPv4 options)
>
> Type of Service and Precedence:
>> By default, copied from the IPv6 Traffic Class (all 8
>> bits). According to [DIFFSERV] the semantics of the
>> bits are identical in IPv4 and IPv6. However, in
>> some IPv4 environments these bits might be used with
>> the old semantics of "Type Of Service and
>> Precedence". An implementation of a translator
>> SHOULD provide the ability to ignore the IPv6 traffic
>> class and always set the IPv4 "TOS" to zero.
>
> Total Length:
>> Payload length value from IPv6 header, plus the size
>> of the IPv4 header.
>
> Identification:
>> All zero.
>
> Flags:
>> The More Fragments flag is set to zero. The Don't
>> Fragments flag is set to one.
>
> Fragment Offset:
>> All zero.

Time to Live:
        Hop Limit value copied from IPv6 header. Since the
        translator is a router, as part of forwarding the
        packet it needs to decrement either the IPv6 Hop
        Limit (before the translation) or the IPv4 TTL (after
        the translation). As part of decrementing the TTL or
        Hop Limit the translator (as any router) needs to
        check for zero and send the ICMPv4 or ICMPv6 "ttl
        exceeded" error.

Protocol:
        Next Header field copied from IPv6 header.

Header Checksum:
        Computed once the IPv4 header has been created.

Source Address:
        If the IPv6 source address is an IPv4-translated
        address then the low-order 32 bits of the IPv6 source
        address is copied to the IPv4 source address.
        Otherwise, the source address is set to 0.0.0.0. The
        use of 0.0.0.0 is to avoid completely dropping e.g.
        ICMPv6 error messages sent by IPv6-only routers which
        makes e.g. traceroute present something for the
        IPv6-only hops.

Destination Address:
        IPv6 packets that are translated have an IPv4-mapped
        destination address. Thus the low-order 32 bits of
        the IPv6 destination address is copied to the IPv4
        destination address.

If any of an IPv6 hop-by-hop options header, destination options
header, or routing header with the Segments Left field equal to zero
are present in the IPv6 packet, they are ignored i.e., there is no
attempt to translate them. However, the Total Length field and the
Protocol field would have to be adjusted to "skip" these extension
headers.

If a routing header with a non-zero Segments Left field is present
then the packet MUST NOT be translated, and an ICMPv6 "parameter
problem/ erroneous header field encountered" (Type 4/Code 0) error
message, with the Pointer field indicating the first byte of the
Segments Left field, SHOULD be returned to the sender.

If the IPv6 packet contains a Fragment header the header fields are
set as above with the following exceptions:

> Total Length:
> > Payload length value from IPv6 header, minus 8 for
> > the Fragment header, plus the size of the IPv4
> > header.
>
> Identification:
> > Copied from the low-order 16-bits in the
> > Identification field in the Fragment header.
>
> Flags:
> > The More Fragments flag is copied from the M flag in
> > the Fragment header. The Don't Fragments flag is set
> > to zero allowing this packet to be fragmented by IPv4
> > routers.
>
> Fragment Offset:
> > Copied from the Fragment Offset field in the Fragment
> > Header.
>
> Protocol:
> > Next Header value copied from Fragment header.

4.2. Translating ICMPv6 Headers into ICMPv4 Headers

All ICMP messages that are to be translated require that the ICMP
checksum field be updated as part of the translation since ICMPv6,
unlike ICMPv4, has a pseudo-header checksum just like UDP and TCP.

In addition all ICMP packets need to have the Type value translated
and for ICMP error messages the included IP header also needs
translation.

The actions needed to translate various ICMPv6 messages are:

> ICMPv6 informational messages:
>
> > Echo Request and Echo Reply (Type 128 and 129)
> > > Adjust the type to 0 and 8, respectively, and adjust the ICMP
> > > checksum both to take the type change into account and to
> > > exclude the ICMPv6 pseudo-header.
> >
> > MLD Multicast Listener Query/Report/Done (Type 130, 131, 132)
> > > Single hop message. Silently drop.

      Neighbor Discover messages (Type 133 through 137)
         Single hop message. Silently drop.

      Unknown informational messages
         Silently drop.

   ICMPv6 error messages:

      Destination Unreachable (Type 1)
         Set the Type field to 3. Translate the code field as
         follows:
             Code 0 (no route to destination):
                 Set Code to 1 (host unreachable).

             Code 1 (communication with destination administratively
             prohibited):
                 Set Code to 10 (communication with destination host
                 administratively prohibited).

             Code 2 (beyond scope of source address):
                 Set Code to 1 (host unreachable). Note that this
                 error is very unlikely since the IPv4-translatable
                 source address is considered to have global scope.

             Code 3 (address unreachable):
                 Set Code to 1 (host unreachable).

             Code 4 (port unreachable):
                 Set Code to 3 (port unreachable).

      Packet Too Big (Type 2)
         Translate to an ICMPv4 Destination Unreachable with code 4.
         The MTU field needs to be adjusted for the difference between
         the IPv4 and IPv6 header sizes taking into account whether or
         not the packet in error includes a Fragment header.

      Time Exceeded (Type 3)
         Set the Type to 11. The Code field is unchanged.

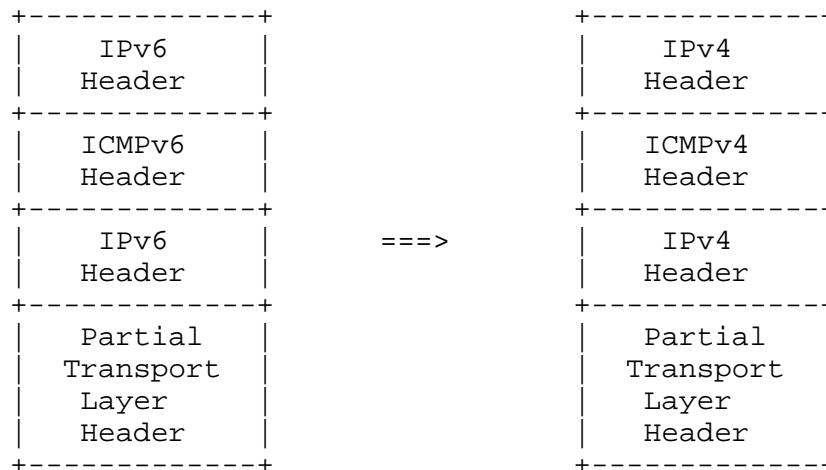   Parameter Problem (Type 4)
         If the Code is 1 translate this to an ICMPv4 protocol
         unreachable (Type 3, Code 2). Otherwise set the Type to 12
         and the Code to zero. The Pointer needs to be updated to
         point to the corresponding field in the translated include IP
         header.

      Unknown error messages
         Silently drop.

4.3. Translating ICMPv6 Error Messages into ICMPv4

   There are some differences between the IPv4 and the IPv6 ICMP error
   message formats as detailed above. In addition, the ICMP error
   messages contain the IP header for the packet in error which needs to
   be translated just like a normal IP header. The translation of this
   "packet in error" is likely to change the length of the datagram thus
   the Total Length field in the outer IPv4 header might need to be
   updated.

```
        +-------------+                    +-------------+
        |    IPv6     |                    |    IPv4     |
        |   Header    |                    |   Header    |
        +-------------+                    +-------------+
        |   ICMPv6    |                    |   ICMPv4    |
        |   Header    |                    |   Header    |
        +-------------+                    +-------------+
        |    IPv6     |        ===>        |    IPv4     |
        |   Header    |                    |   Header    |
        +-------------+                    +-------------+
        |   Partial   |                    |   Partial   |
        |  Transport  |                    |  Transport  |
        |    Layer    |                    |    Layer    |
        |   Header    |                    |   Header    |
        +-------------+                    +-------------+
```

            IPv6-to-IPv4 ICMP Error Translation

   The translation of the inner IP header can be done by recursively
   invoking the function that translated the outer IP headers.

4.4. Knowing when to Translate

   When the translator receives an IPv6 packet with an IPv4-mapped
   destination address the packet will be translated to IPv4.

5. Implications for IPv6-Only Nodes

   An IPv6-only node which works through SIIT translators need some
   modifications beyond a normal IPv6-only node.

   As specified in Section 1.3 the application protocols need to handle
   operation on a dual stack node. In addition the protocol stack needs
   to be able to:

o Determine when an IPv4-translatable address needs to be allocated
   and the allocation needs to be refreshed/renewed. This can
   presumably be done without involving the applications by e.g.
   handling this under the socket API. For instance, when the
   connect or sendto socket calls are invoked they could check if the
   destination is an IPv4-mapped address and in that case
   allocate/refresh the IPv4-translatable address.

o Ensure, as part of the source address selection mechanism, that
   when the destination address is an IPv4-mapped address the source
   address MUST be an IPv4-translatable address. And an IPv4-
   translatable address MUST NOT be used with other forms of IPv6
   destination addresses.

o Should the peer have AAAA/A6 address records the application (or
   resolver) SHOULD never fall back to looking for A address records
   even if communication fails using the available AAAA/A6 records.
   The reason for this restriction is to prevent traffic between two
   IPv6 nodes (which AAAA/A6 records in the DNS) from accidentally
   going through SIIT translators twice; from IPv6 to IPv4 and to
   IPv6 again. It is considered preferable to instead signal a
   failure to communicate to the application.

6. Security Considerations

   The use of stateless IP/ICMP translators does not introduce any new
   security issues beyond the security issues that are already present
   in the IPv4 and IPv6 protocols and in the routing protocols which are
   used to make the packets reach the translator.

   As the Authentication Header [IPv6-AUTH] is specified to include the
   IPv4 Identification field and the translating function not being able
   to always preserve the Identification field, it is not possible for
   an IPv6 endpoint to compute AH on received packets that have been
   translated from IPv4 packets. Thus AH does not work through a
   translator.

   Packets with ESP can be translated since ESP does not depend on
   header fields prior to the ESP header. Note that ESP transport mode
   is easier to handle than ESP tunnel mode; in order to use ESP tunnel
   mode the IPv6 node needs to be able to generate an inner IPv4 header
   when transmitting packets and remove such an IPv4 header when
   receiving packets.

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